



A Recipe for Cryptographic TEEs...

- 1. Eliminate *all* software vulnerabilities
- But all software is (eventually) hackable!
- Approach: *no S/W in the enclave*
- 2. Silence side channels
- Control, memory, timing, and uArch
- Approach: provably side-channel free enclave
- 3. Encrypt sensitive data everywhere else
 - Eliminates trust for all remaining S/W and H/W • Approach: encrypt on exit of enclave





Use Case: Medical Data Sharing

- Privacy-preserving smartwatchbased heart monitoring
- **Data encryption:** biometric data is encrypted at the sensor, sent to server encrypted
- *Encrypted computation:* analysis of biometric data is performed on encrypted data, without giving access to the server or its operators
- · Guardrailing: attackers cannot manipulate analysis algorithms to trigger false warnings
- Safe datagrants: arrhythmia analysis algorithm can expose potential health warnings to permit notification/text message to the user



Crypto-strength defenses: Not vulnerable to S/W hacking or digital side channels!

TrustForge: A Cryptographically Secure Enclave

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This 190k-gate functional unit implements always-encrypted computation, nullifying hacking and advancing privacy technology.

TrustForge: Inside the Zone of Trust

- TrustForge enclave exports a public key, PKI-encrypted keys are decrypted in enclave
- Enclave supports RISC-like operations that operate directly on encrypted data
- Enclave is decrypting, processing, checking, and re-encrypting secret data *without S/W* vulnerabilities or digital side channels
- Enclave is protected from physical attacks, and *data is encrypted everywhere else*





Slowdown	10000000	
	1000000	
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	1	

Programming for the TrustForge Enclave

- TrustForge extends development language with encrypted variables
- *Bit-for-bit compatible*, but encrypted
- Support for *integers, floating-point*, Booleans, and strings
- Encrypted ops return encrypted results
- Decision processing on encrypted variables
 implemented with *CMOV primitive*
- Secret-dependent array indexing implemented with **ORAM primitives**

<u>}</u>	Type Class	C++ Data Types		
	Integer	enc_int, enc_uint, enc_long, enc_ulong		
	Floating Point	enc_float, enc_double		
	Boolean	enc_bool		
	String/Char	enc_char, enc_char, enc_string		
	$\begin{array}{cccc} 1 & \mathbf{x} &= & \operatorname{enc}_{-} \mathbf{c} \\ 2 & \mathbf{y} &= & \operatorname{enc}_{-} \mathbf{c} \end{array}$	<pre>cmov(secret, x+1, x); cmov(!secret, y+1, y);</pre>		
for	<pre>(int i=0; i<1 rot = onc cmout)</pre>	<pre>.en(arr); i++;) (i==secret arr[i] ret)</pre>		

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StForge Security Analysis

 Worked with In-Q-Tel & Blue Bastion. Findings Summary Blue Bastion pen tested for 3 months A total of <mark>zero</mark> (0) findings were identified in this repor Low Information Red team had full access to all IP Zero vulnerabilities found Formal verification with Princeton Security Verification of Low-Trust Architectures • Secure for *any program* on a ABSTRACT Low-trust architectures work on, from the viewpoint of softwa ways-encrypted data, and significantly reduce the amount of ha specific implementation re trust to a small software-free enclave component. In this pa e perform a complete formal verification of a specific low-trust we perform a compare format verification of a specific tow that architecture, to show that the design is secure against direct data disclosures and digital side channels for all possible programs. We first define the secu-rity requirements of the ISA of SE low-trust architecture. Looking ative instructions, termed the native ISA, and a set of secure in structions termed the SE ISA. The native ISA contains insecure Zero vulnerabilities found pwards, this ISA serves as an abstraction of the hardware for the Forware, and is used to show how any program comprising these structions cannot leak information, including through digital side hannels. Looking downwards this ISA is a specification for the ardware, and is used to define the proof obligations for any RTL • Proofs to appear in ACM CCS 2023 is the design of this SE Enclave and ISA extension that ensures t SE Background – Section 2 Threat Model – Section 3 implementation arising from the ISA-level security requirements. These cover both functional and digital side-channel leakage. Next,